

# Typechecking of Overloading

## in Programming Languages and Mechanized Mathematics

Arthur Charguéraud and Martin Bodin  
with Louis Riboulet and Jana Dunfield

Inria

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## Contribution

Present the first overloading resolution algorithm accompanied with a polynomial complexity bound.

Bound in terms of the size of the output typed tree.

Algorithm complete with respect to *iteratively solvable programs*, i.e. solvable with a chain of deduction steps (otherwise problem is NP-hard).

Assumes annotation on polymorphic definitions.

# Contribution

Implementation in OCaml, parses extended OCaml syntax.

Overloading of functions, constants, constructors and record fields.

Supports instances with assumptions, e.g., instance for `show on 'a list` requires instance of `show` for `'a`.

Output fully typed AST, with symbols resolved.

Extracts to standard, executable OCaml code.

# Algorithm overview

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2. Collect symbols in traversal order, overloaded functions both ways
3. Iteratively try to resolve symbols, in a carefully chosen order.

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## Resolution rule:

Assume symbol  $x$  with a type  $T$  constrained by the context.

Assume  $x$  has candidate instances  $(v_1 : T_1), \dots, (v_n : T_n)$ .

Then  $x$  resolves to  $v_i$  if  $T$  unifies with  $T_i$  but not with  $T_j$  for  $j \neq i$ .

## No backtracking

If an instance has *assumptions*, they are investigated only after the choice of the instance is made.

# Overlapping

An instance can be declared to *override* a prior instance: in case of overlap, the second will be preferred.

Ex: generic instance for show-list, but specific instance for show-list-of-char, to print using string notation.

## Possible outcomes

Termination is ensured by limiting the dependency depth.

- ▶ Success: all resolved
- ▶ Ill-typed: type conflict, or no remaining candidate
- ▶ Stuck: cannot make any deduction.
- ▶ MaxDepthExceeded: unexpectedly high search depth.

# Soundness

*Typechecking*: ML-typechecking followed by iterated symbol resolution.

*Extracted program*: the program obtained by replacing overloaded symbols with the values they resolved to.

## Theorem (Type soundness)

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*If a program successfully typechecks, then the extracted program is well-typed in ML.*

## Theorem (Non-ambiguity)

*If our typechecking algorithm succeeds, then no other instantiation of the overloaded symbols leads to a well-typed ML program.*

## Implication of success and failure

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*If our algorithm successfully typechecks a program, then it finds the only instantiation of the overloaded symbols that leads to a well-typed ML program. In particular, the program is solvable.*

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### Theorem (Implication of a type error)

*If our typechecking algorithm raises an error, then no instantiation of the overloaded symbols leads to a well-typed ML program. In particular, the program is not solvable.*

# Completeness

*Iteratively solvable program*: iff there exists an ordering  $x_1, x_2, \dots, x_n$  of the overloaded symbols such that, after resolving the  $x_{k < i}$  it is possible to resolve  $x_i$  without inspecting the candidates for the  $x_{k > i}$ .

*Assumption depth*: maximal number of nesting of assumption in the resolution of a solvable program.

## Theorem (Characterization of successful typechecking)

*Our algorithm with parameter max-depth set to  $D$  successfully typechecks a program if and only if this program is iteratively solvable and with an assumption depth not exceeding  $D$ .*

## Details

Ask for our soon-to-be-submitted article for details.