Rewriting under binders, comfortably

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$$\sum_{i=0}^{n} i = \sum_{i=0}^{n} \sqrt{i^2}$$

- Math teacher proofs are very different,
- A proposed solution.

The context

- Mathematical constructions like integrals and iterated sums have bound variables
- From the formal point of view, a bound variable does not really exist
- Type theory promotes *Leibniz* equality as the main tool to reason modulo equality

especially for rewriting

But Leibniz equality requires objects that really exist

Discrepancy in idioms

$$\sum_{i=0}^{n} i = \sum_{i=0}^{n} \sqrt{i^2}$$

The math teacher's proof (I believe)

• Replace $\sqrt{i^2}$ with *i* in the right-hand side sum.

Note that the sum ranges over positive values

The formally verified proof

1. Establish
$$\forall i, 0 \leq i \leq n \Rightarrow i = \sqrt{i^2}$$

- 2. For this, fix *i* such that $0 \le i \le n$,
- 3. Then $i = \sqrt{i^2}$ (by some proof),

4. then apply the extensionality lemma for sums:

$$\forall fg, (\forall i, 0 \leq i \leq n \Rightarrow f(i) = g(i)) \Rightarrow \sum_{i=0}^{n} f(i) = \sum_{i=0}^{n} g(i)$$

The curse of α -conversion

• There is no doubt that, if *i* exists and is larger than 0, $i = \sqrt{i^2}$,

Leibniz says: if n = m, you can replace n with m in any formula

• But the numbers *i* and \sqrt{i} do not even exist in the formula

$$\sum_{i=0}^{n} \sqrt{i^2}$$

Bound variable names do not count for logical reasoning

$$\sum_{i=0}^n \sqrt{i^2} = \sum_{j=0}^n \sqrt{j^2}$$



A preliminary solution

- Make the sentence Replace \sqrt{i^2} with i in the sum. understandable by the proof system
- **b** Do not work modulo α -conversion
- 1. Recognize that $\sqrt{i^2}$ is not well-formed because we are missing a variable with the name i
- 2. By scanning the formula, detect that *i* is bound in at least one place,
- 3. Search for instances of $\sqrt{i^2}$ in the multiple places where this may occur
- 4. Do this again if there are nested binding patterns
- 5. Every time one enters inside an operator with bound variables, apply a suitable "extensionality" theorem

Example using the solution

DEMO

dépot git, fichier d'expérience

A prototype implementation

 Required an extension of the Elpi meta-programming language

Authorize passing "open terms" as argument to tactics

- An open term is well typed in an extension of the context
- ► Example, if *i* does not exist in the context √*i*² is not well-typed, but λ*i* : ℝ, √*i*² is well typed
- The tactic receives two open terms, which can be viewed has a rewrite rule
- the context is search for a subcontext where:
 - All "open variables" are accounded
 - The left-hand side of the rewrite rule occurs

Building a proof

Default extensionality: two functions that are equal everywhere can be substituted for each other
Axiom functional_extensionality provided by Rocq
Ad hoc extensionality: compare functions only on a subset
For integrals: the subset is the interval between the bounds
For discrete sums with integer bounds: the subset is the intersection of the integers and the interval between the bounds

Future work

- Provide a comfortable interface to add new ad-hoc extensionality principles
- Rely on setoid rewrite, advanced location selection
- Handle goals that are not equalities
- perform replacement modulo orders
- Find links with observational type theory